

An Introduction to Synchronous Data Flow Model

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CPEG-852: Advanced Topics in Computing Systems

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Outline

- Synchronous Data Flow Model
 - Definition
 - Example
- Periodic Schedule and Consistency
- Stream Programming Language
 - Structured SDF
- Apply SDF to Bigdata

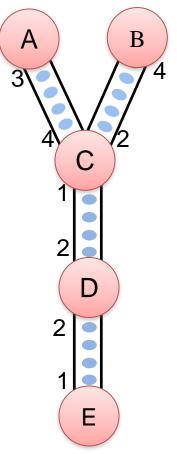


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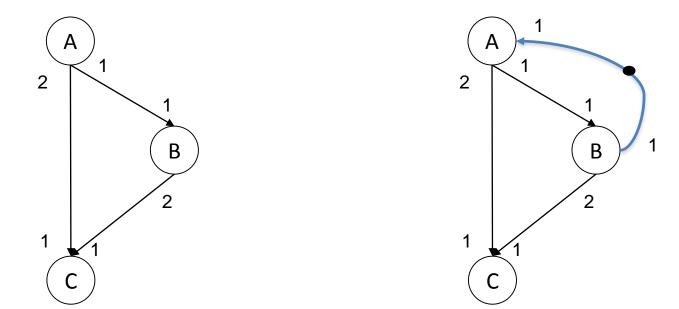
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- Synchronous Data Flow (SDF) is represented as a graph
 - Node (actor): Computation
 - Edge: First In First Out (FIFO) Queue
- Each edge has two weights: produce rate and consume rate
- Each edge can also have initial data
- Formal: Tuple<N, E, E_{p,c,i}>,
 - N: node
 - E: edge
 - E_{p,c,i}: Produce rate, consume rate and initial data



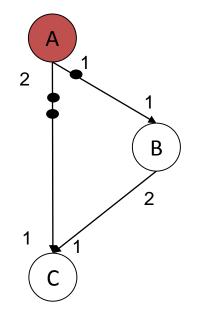




SDF with no initial tokens

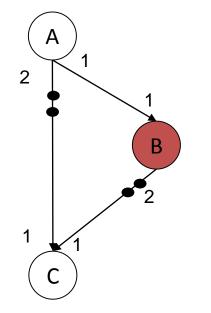
SDF with initial token and loop





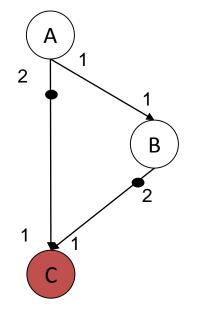
A firing





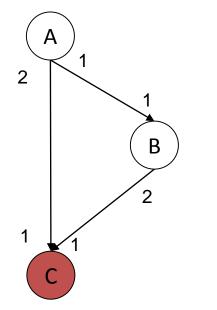
A firing, B firing





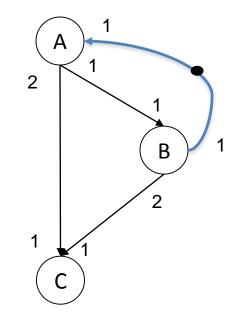
A firing, B firing, C firing





A firing, B firing, C firing, C firing

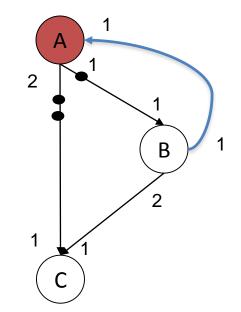




SDF with initial token and loop

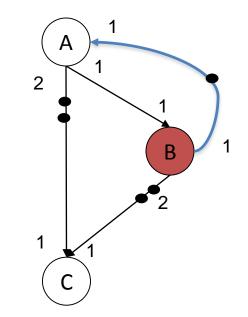


Synchronous Data Flow Model



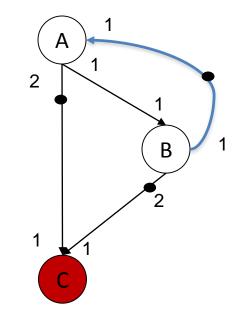
A firing





A firing, B firing

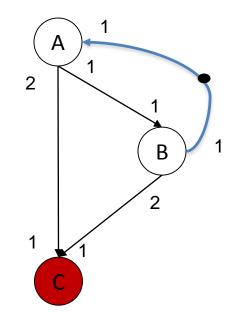




A firing, B firing, C firing



Synchronous Data Flow Model



A firing, B firing, C firing, C firing



- Question: Can any SDF graph find a firing sequences that makes the state of the graph no changed?
 - State of the graph means: the tokens on each edge are clean, no more no less.
 - Which leads to SDF Consistency Problem



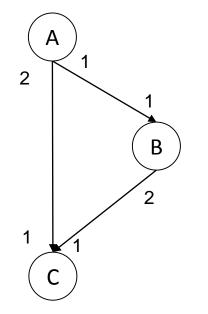
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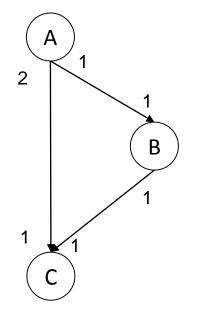
- Firing sequence of a SDF is called a *schedule*
- A *periodic schedule* of an SDF clears all channels and return to its initial status after each node repeats execution a specified finite number of times
- Periodic schedule, permit SDF can process unbounded data with bounded memory
- A SDF is Consistent if a periodic schedule exists





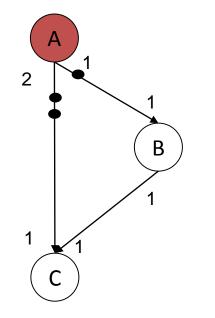
Periodic Schedule: ABCC AB2C





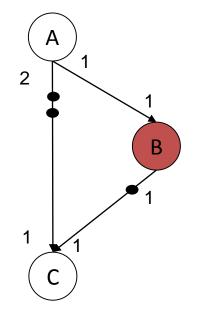
Can you find the periodic schedule?





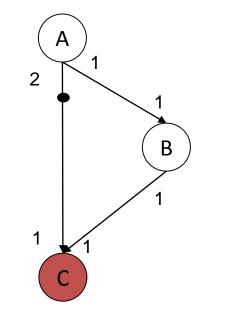
А





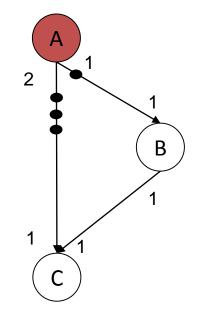
A, B





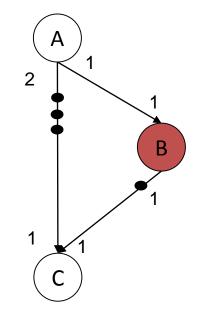
A, B, C





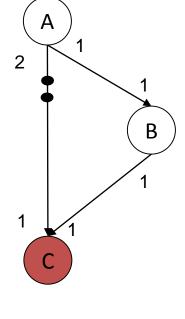
A, B, C, A





A, B, C, A, B



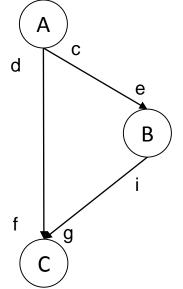


Tokens in channel (A-C) is accumulating which makes the channel unbounded

A, B, C, A, B, C

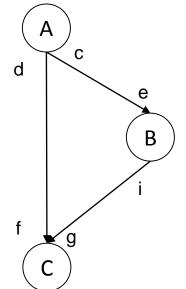
Inconsistent!





Problem: Given a general SDF, how can we know it has periodic schedule or not?



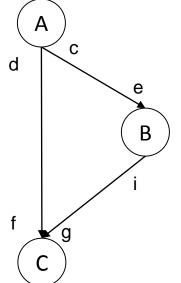


Topology Matrix

- Each row presents the edge
- Each column presents a node
 - (i, j): the number of data items placed on i after each invocation of j
 - If i is an input channel for j, element (i, j) is negative

$$\begin{array}{ccc} A & B & C \\ c & -e & 0 \\ d & 0 & -f \\ 0 & i & -g \end{array} \right) \begin{array}{c} A \rightarrow B \\ A \rightarrow B \\ A \rightarrow C \\ B \rightarrow C \end{array}$$

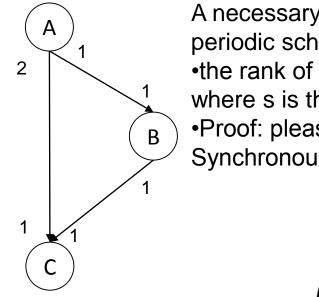




A necessary condition for the existence of a periodic schedule •the rank of the topology matrix is s – 1, where s is the number of nodes •Proof: please refer to "Lee's 87 paper: Synchronous Data Flow"

$$\begin{pmatrix} c & -e & 0 \\ d & 0 & -f \\ 0 & i & -g \end{pmatrix} A \rightarrow B A \rightarrow B A \rightarrow C B \rightarrow C B \rightarrow C$$





A necessary condition for the existence of a periodic schedule •the rank of the topology matrix is s – 1, where s is the number of nodes •Proof: please refer to "Lee's 87 paper: Synchronous Data Flow"

$$\begin{array}{ccc} 1 & -1 & 0 \\ 2 & 0 & -1 \\ 0 & 1 & -1 \end{array} \begin{array}{c} A \rightarrow B \\ A \rightarrow C \\ B \rightarrow C \end{array}$$

Rank=3 > 2

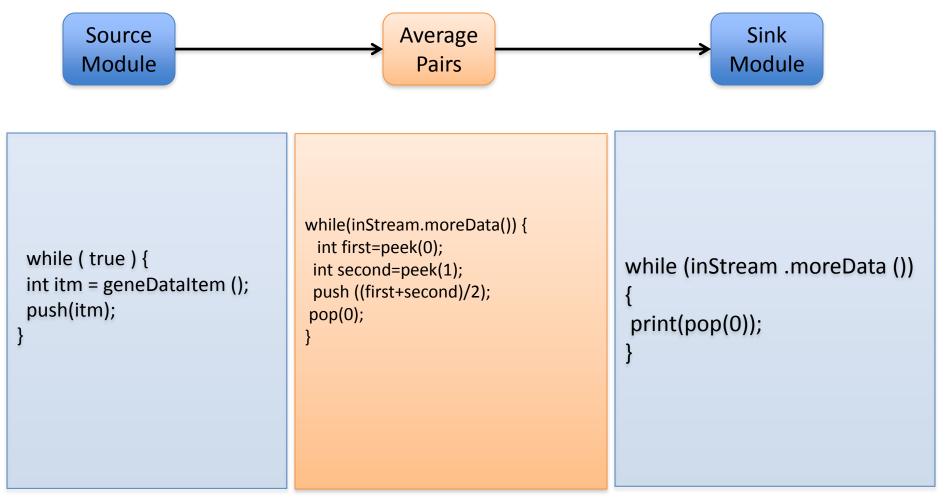


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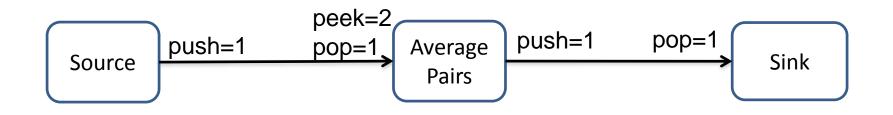
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An Stream Example—Average Pairs







Extend the SDF to support "peek" sematic























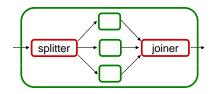
Structured SDF In StreamIt

• Filter

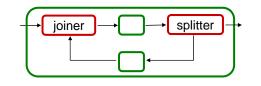
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• Split-Join

Pipeline

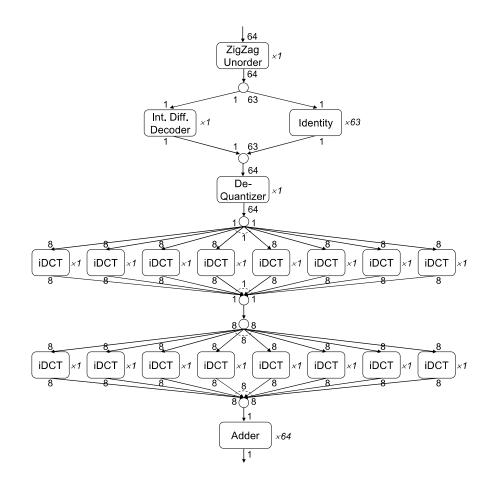


Feedback Loop





Part of JPEG transcoder



Courtesy by: PACT 2010 paper of streamit: "An Empirical Characterization of Stream Programs and its Implications for Language and Compiler Design "



Weakness of SDF

- Does not support condition (branch)
- Does not support recursion—because it is a static dataflow model
- But still the model is used widely in many application fields



Some Projects Based on SDF model

- Early Ptolemy Project at UC Berkeley
 - Software Synthesis for Embedded system
- StreamIt at MIT
 - streaming program language and compiler
- InforStream and SPL
 - IBM streaming computing product
- Our work COStream
 - hierarchical date flow programming language and compiler
- OpenStream
 - language and compiler support for streaming in OpenMP



Homework

- Write a Fibonacci number generator using Synchronous Data Flow Model
 - Pseudo code for each node in SDF using "peek, push and pop" statements
 - Push Token: PPT animation to show how the tokens flow in SDF graph
 - Periodic Schedule of the SDF



Reference

- [1] Early Ptolemy Project at UC Berkeley
 - <u>http://ptolemy.eecs.berkeley.edu/projects/index.htm</u>
- [2] StreamIt at MIT
 - <u>http://groups.csail.mit.edu/cag/streamit/</u>
- [3] InforStream and SPL
 - <u>http://www-03.ibm.com/software/products/en/infosphere-streams</u>
- [4] COStream
 - <u>http://www.capsl.udel.edu/pub/doc/papers/dfm12.pdf</u>
- [5] OpenStream
 - <u>http://openstream.info/</u>